

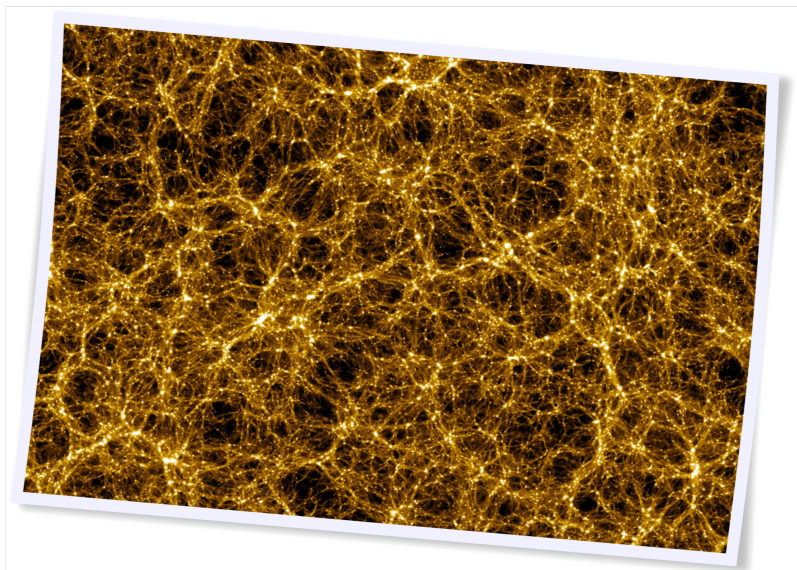
The Single Parameter Fast Multipole Method

A Coulomb solver for the masses.

September 4, 2012 | Ivo Kabadshow | Holger Dachselt |

Real World Problems

Astrophysics: Structure Formation



Real World Problems

Astrophysics: Galaxy Formation



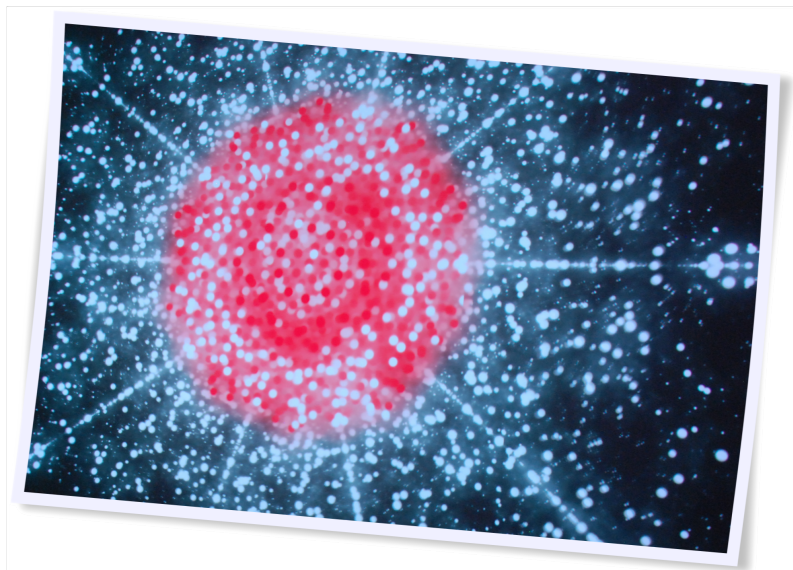
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Real World Problems

Laser-Plasma Physics: Coulomb Explosion



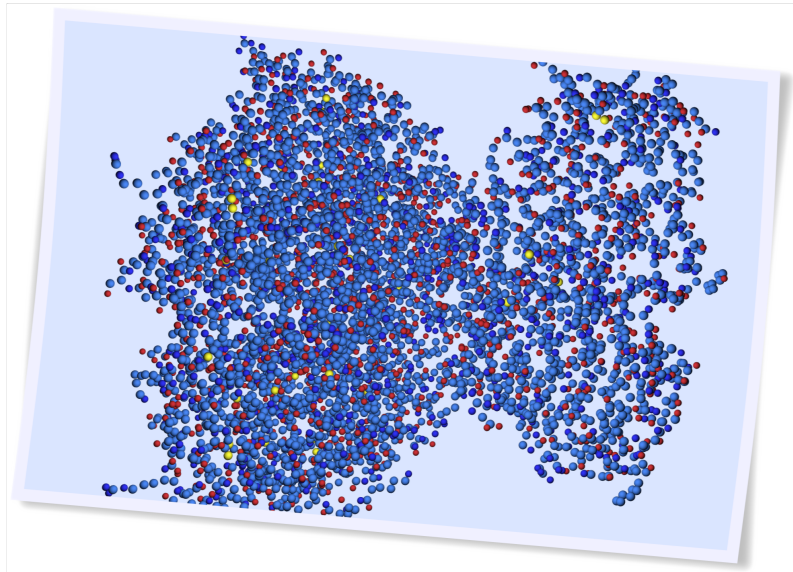
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Real World Problems

Softmatter Physics: Biological Systems



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The Coulomb Problem

Classical Force Calculation

Direct Summation

$\mathcal{O}(N^2)$

$$\mathbf{F}(\mathbf{r}_i) = q_i \sum_{j=1}^N \frac{q_j}{r_{ij}^3} \mathbf{r}_{ij}$$

In a Simulation

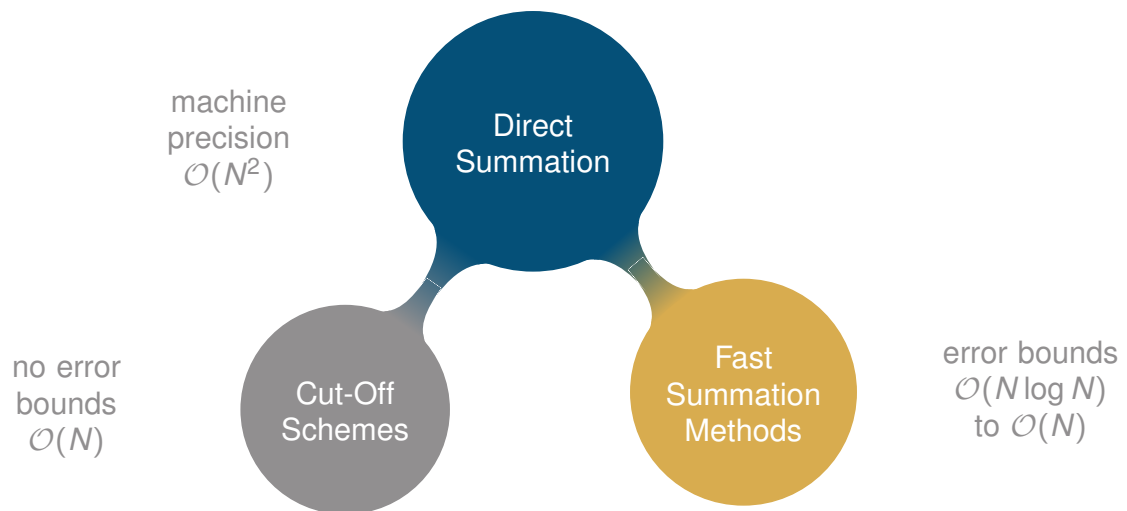
- 95% – 99% of time spend in force calculation
- 1% – 5% of time spend in integration and post processing

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Alternatives To Direct Summation

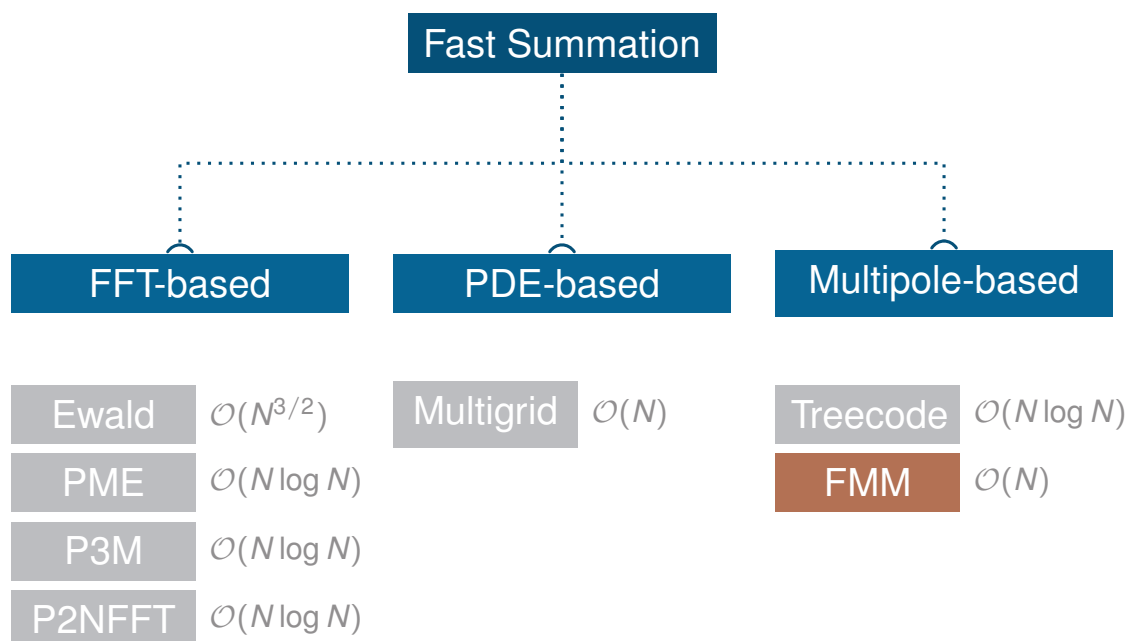


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Classification of Fast Summation Methods



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The Coulomb Problem

Classical Scheme

Direct Summation

$\mathcal{O}(N^2)$

$$E_C = \sum_{i=1}^{N-1} \sum_{j=i+1}^N \frac{q_i q_j}{r_{ij}}$$

Fast Summation Algorithms

Fast Multipole Method

$\mathcal{O}(N)$

$$E_C = \sum_{\text{level}} \sum_A \sum_{B(A)}^p \sum_{l=0}^l \sum_{m=-l}^l \sum_{j=0}^p \sum_{k=-j}^j (-1)^j \omega_{lm}^A(\mathbf{a}) M2L(\mathbf{R}) \omega_{jk}^B(\mathbf{b})$$

Idea: factorization of inverse distance

The Fast Multipole Method

There's no such thing as a free lunch

Pro



- Complexity reduces to $\mathcal{O}(N)$
- Computation time $t = f(\epsilon)$ is a function of requested precision.

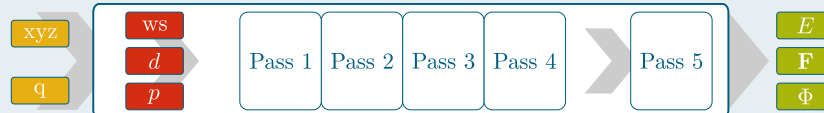
Con



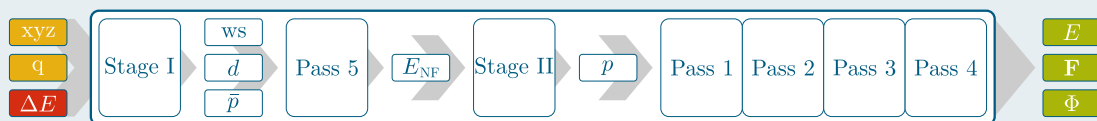
- Three new parameters are introduced
 - Tree depth (d)
 - Separation Criterion (ws)
 - Number of poles (p)
- Precision dependency $\epsilon = g(ws, d, p)$ is not known explicitly

Towards a Single-Parameter FMM

Original FMM Workflow



Enhanced FMM Workflow with Error Control



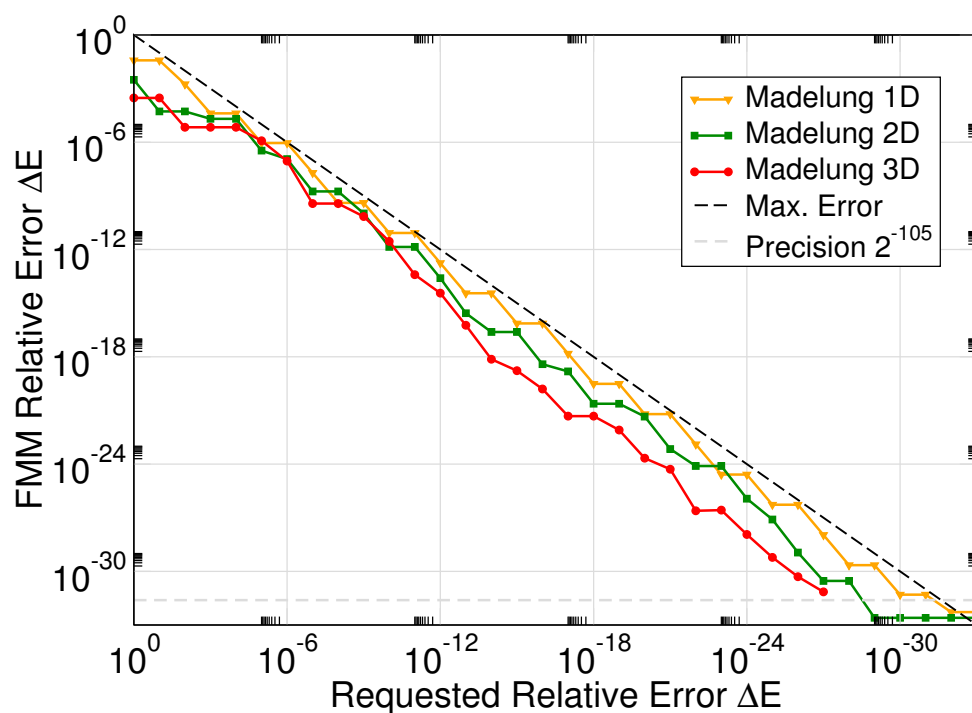
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Quadruple Precision Check

1D, 2D and 3D Periodicity

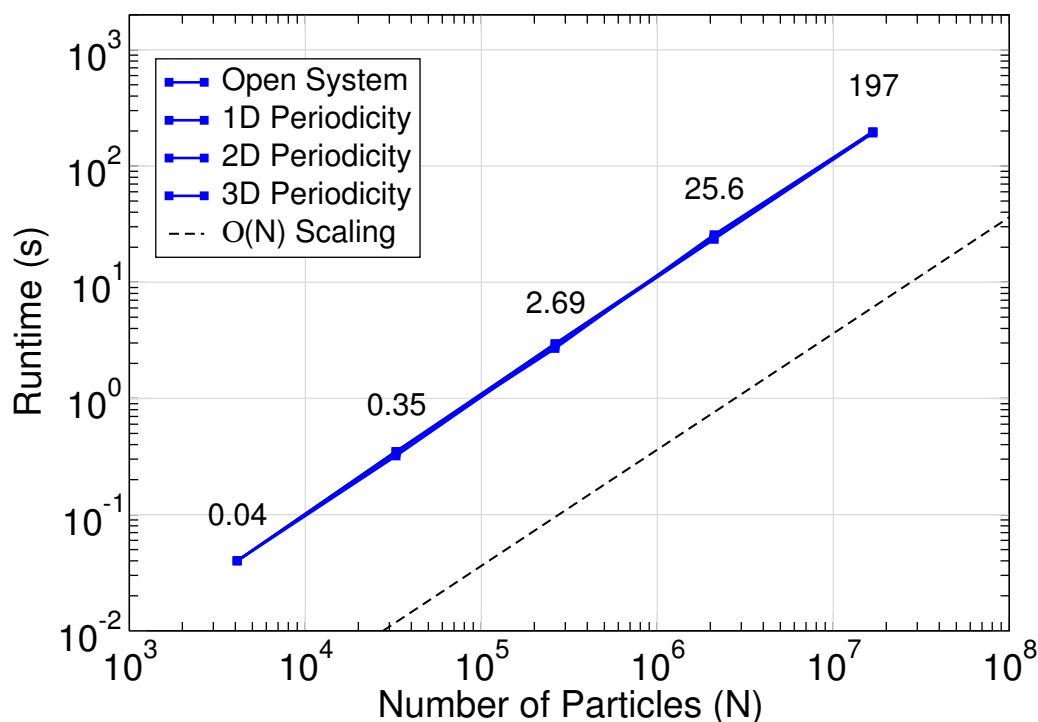


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Scaling with N , $\Delta E_{\text{rel}} = 10^{-5}$



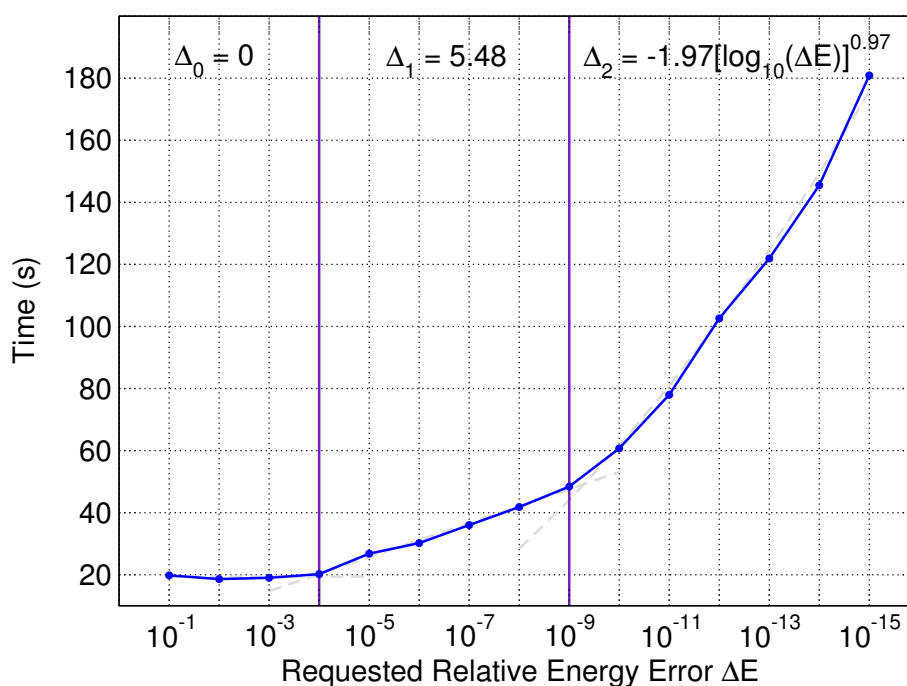
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Precision Scaling

2.097.152 Particles - Total Computation Time



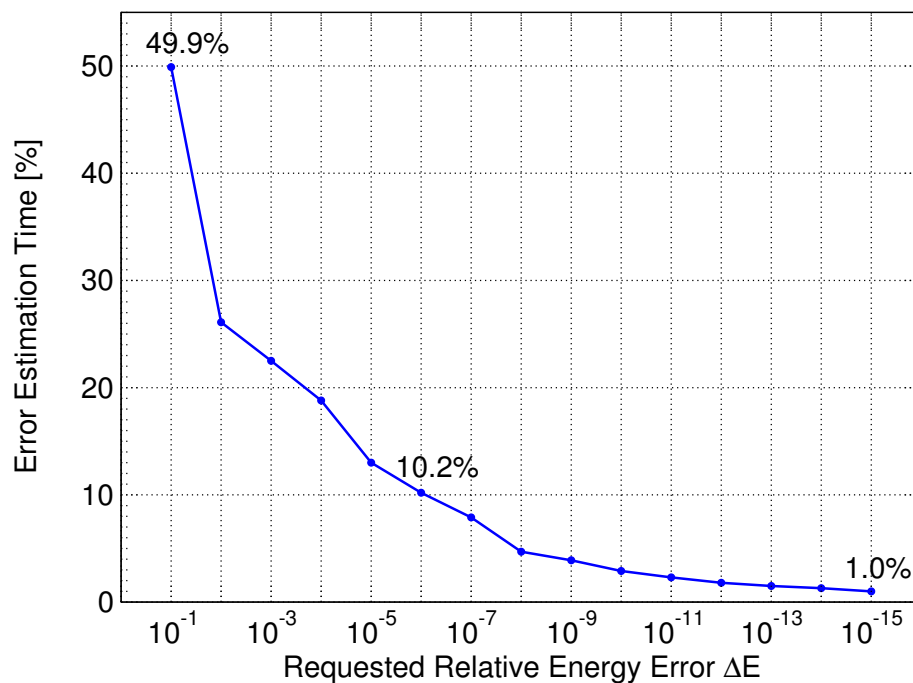
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Tuning Costs

2.097.152 Particles - Error Estimation Scheme Runtime



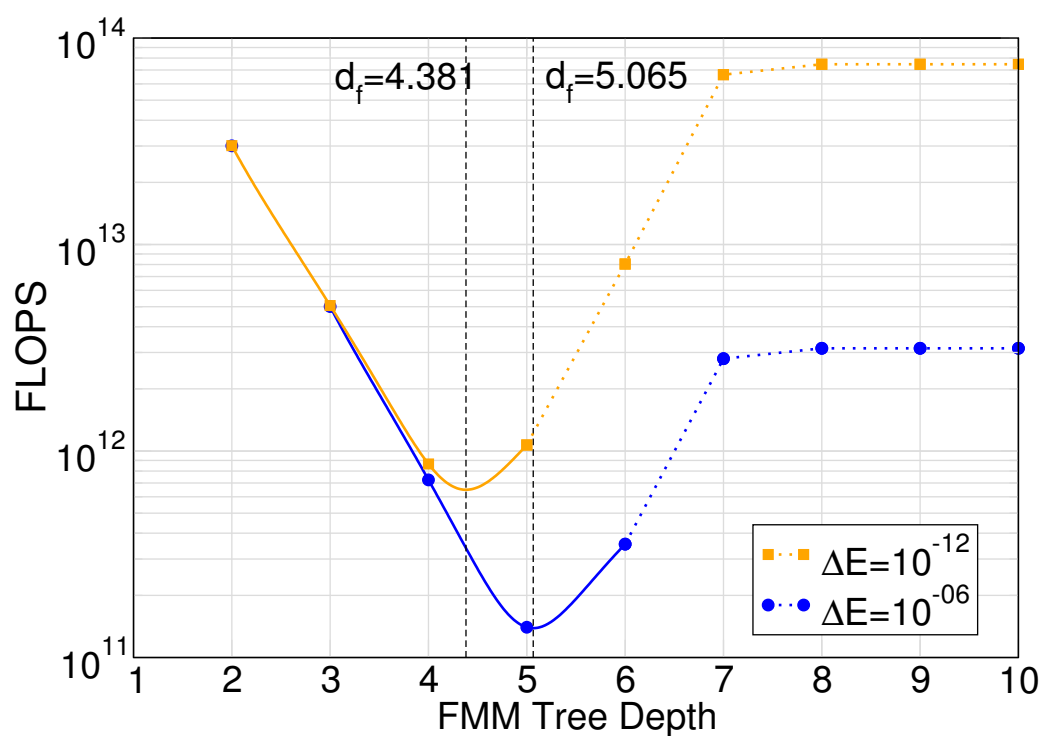
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FMM Runtime Optimization

Trade DIV and SQRT for ADD and MULT



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Going Parallel

Limiting Details

- Fast summation methods trade memory for speed
 - Additional memory again limits system size
 - Critical on distributed memory systems like BG/P

Improve Weak Scaling

- Reduce algorithmic overhead
 - Don't provide storage for empty regions of space
 - Don't store redundant data (neighbor lists etc.)
- Reduce parallelization overhead
 - Latency: overlap communication with computation
 - Memory: avoid $\mathcal{O}(p)$ data structures

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Parallelization Properties

Use ...

- data locality
- fastest one-sided memory access (put) via OSPRI/ARMCI
- minimal communication functionality
- distribute load via space-filling curve
- MPI (inplace) for collectives only (allgather, allreduce, barrier)

Things to Avoid

- Avoid domain decomposition to distribute work
- Avoid accumulate and non-contiguous send operations
- Avoid global counters

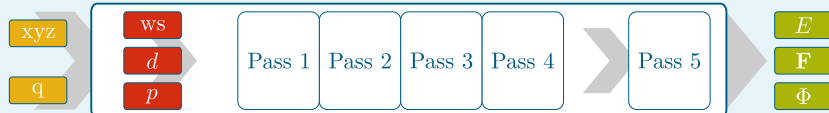
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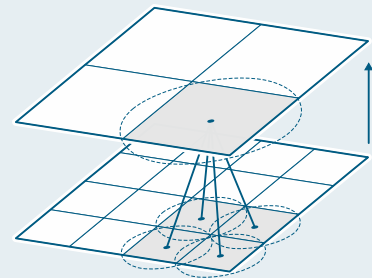
FMM Workflow

FMM Workflow



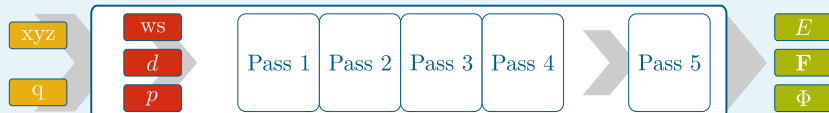
Pass 1

- Setup FMM tree and expand multipoles
- Shift multipoles to root node (M2M)



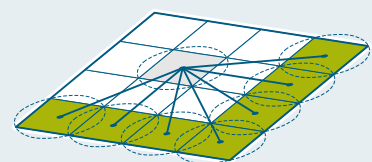
FMM Workflow

FMM Workflow



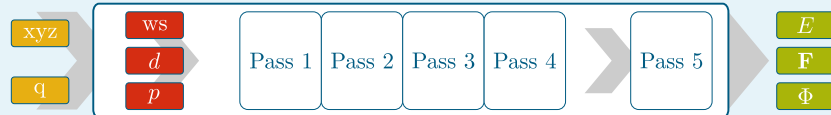
Pass 2

- Translate multipole moments into Taylor coefficients (M2L) at each level of the tree



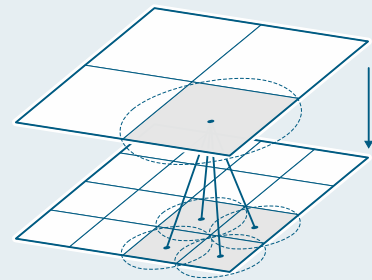
FMM Workflow

FMM Workflow



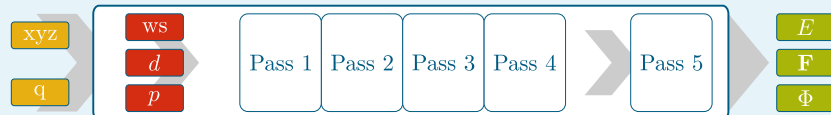
Pass 3

- Shift Taylor coefficients to the leaf nodes (L2L)



FMM Workflow

FMM Workflow

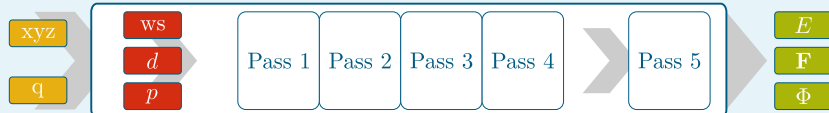


Pass 4

- Computation of the far field energy, forces, potentials

FMM Workflow

FMM Workflow



Pass 5

- Computation of the near field energy, forces, potentials

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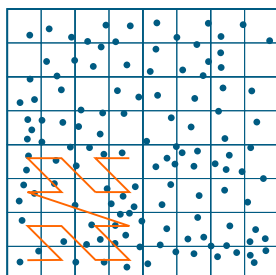
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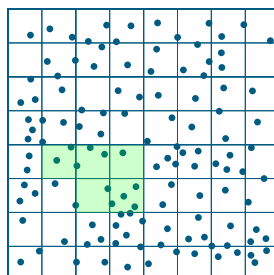
Data Layout

Different Local Domains

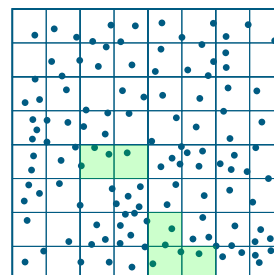
- All boxes are sorted along a space-filling curve
- Static load balancing is essential



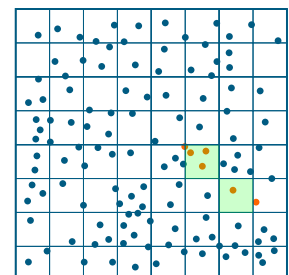
All boxes are
z-ordered



Compact
Domain



Two Domains



Splitted parti-
cles/multipoles

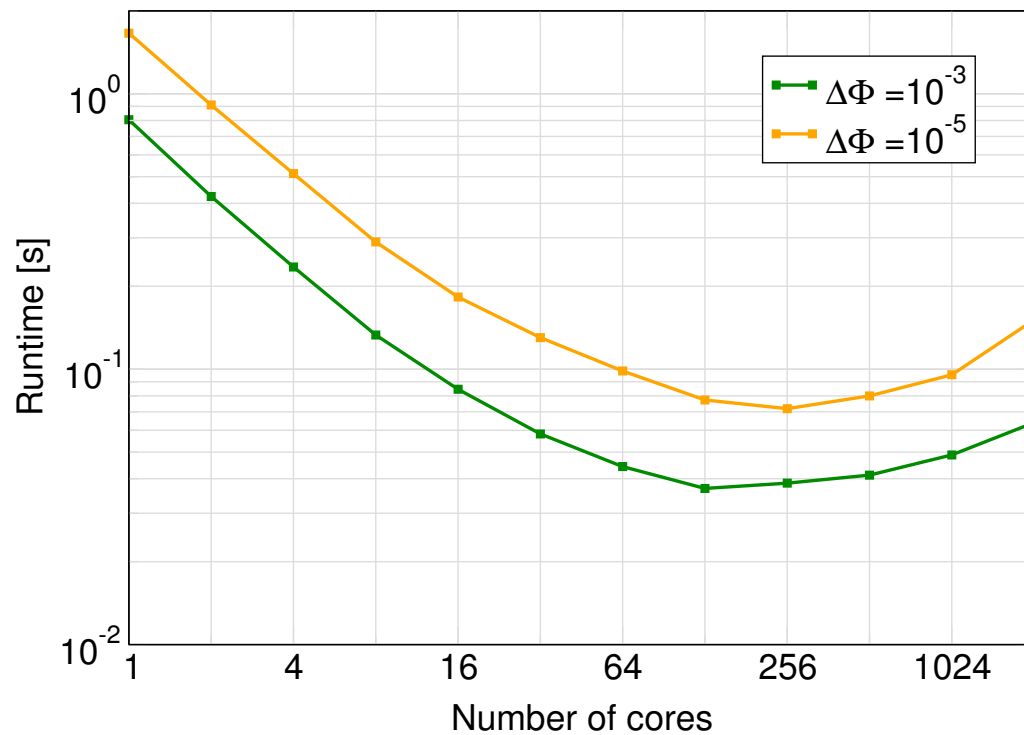
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Strong Scaling On Jugene

FMM with 3D Periodic Boundaries, 8.100 Particles



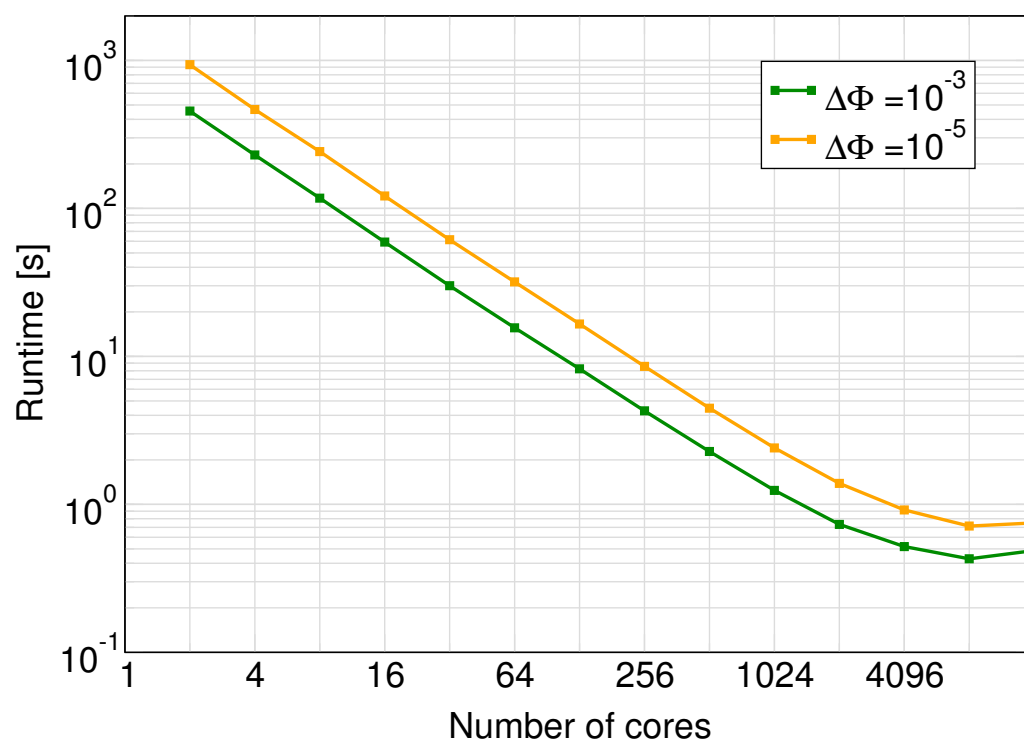
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Strong Scaling On Jugene

FMM with 3D Periodic Boundaries, 9.830.400 Particles



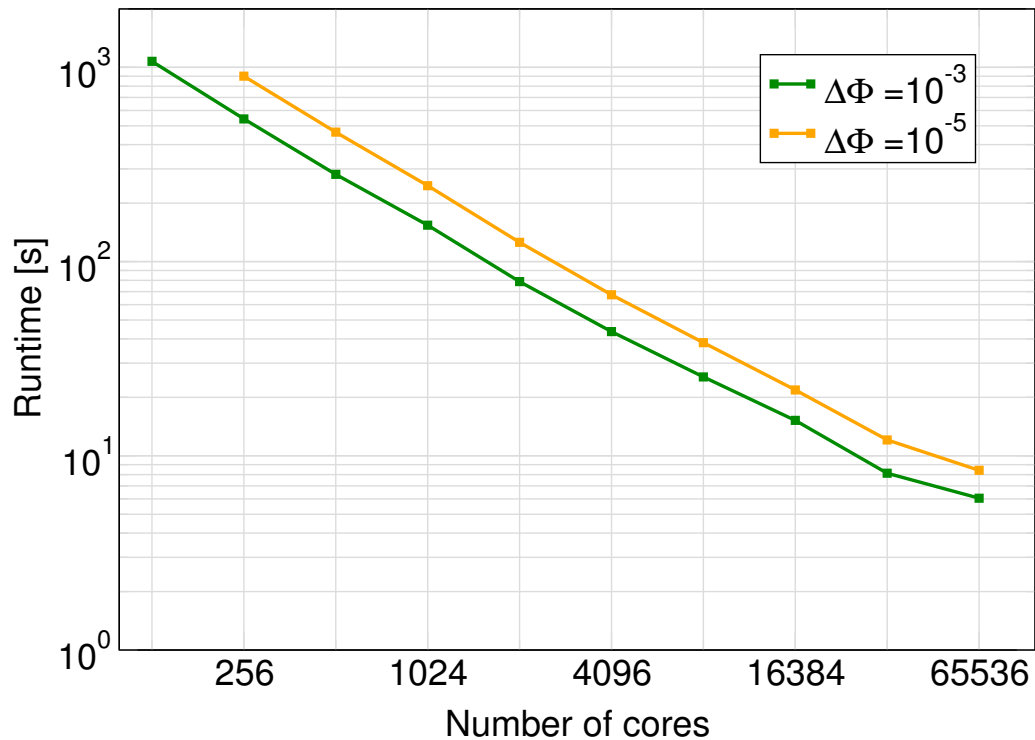
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Strong Scaling On Jugene

FMM with 3D Periodic Boundaries, 1.012.500.000 Particles



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Lower Memory Bound

How many memory would a direct calculation need?

Input (Floats) Size

- Coordinates $xyz(3, n)$ 12 Bytes
- Charges/Masses $q(n)$ 4 Bytes

Output (Floats) Size

- Gradient/Forces/Field $grad(3, n)$ 12 Bytes

Total Memory per Particle Size

28 Bytes

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How Many Particle Can Fit On The Full BGP?

Upper Bound With Direct Summation

$\mathcal{O}(N^2)$

- Available memory $\approx 132\text{ TB}$ (472MB per core)
- 5.2 trillion ($5.2 \cdot 10^{12}$) particles would fit
→ 32.000 years to compute all forces

Direct Summation Introduces Errors

- Machine epsilon $\epsilon = 5.96 \times 10^{-8}$ (single precision)
- Round-off error accumulation $\mathcal{O}(\sqrt{\epsilon N^2})$
→ No significant figure will remain

Pushing The Limits Beyond 1.000.000.000.000

Passing The One-Trillion Particle Limit

System Characteristics

- 1.030.607.060.301 particles
- 32768 nodes (SMP mode, 32768 cores)

Results

- 3285s runtime, unsorted data (9576 particles/second/core)
- 2203s runtime, presorted data (14276 particles/second/core)

Passing The Two-Trillion Particle Limit

System Characteristics

- 2.010.394.559.061 particles
- 73728 nodes (VN mode, 294912 cores)

Results

- 2288s runtime, unsorted data (2279 particles/second/core)
- 530s runtime, presorted data (12862 particles/second/core)

Passing The Three-Trillion Particle Limit

System Characteristics

- 3.011.561.968.121 particles
- 73728 nodes (VN mode, 294912 cores)
- $0.254 \cdot 10^{27}$ FLOPS direct

Results

- (38.00) 44.16 Bytes/particle (16.2% parallelization overhead)
- 3812s runtime, unsorted data (2755 particles/second/core)
- 715s runtime, presorted data (14687 particles/second/core)
 ≈ 4.3 billion particles per second
- $0.410 \cdot 10^{17}$ FLOPS via FMM

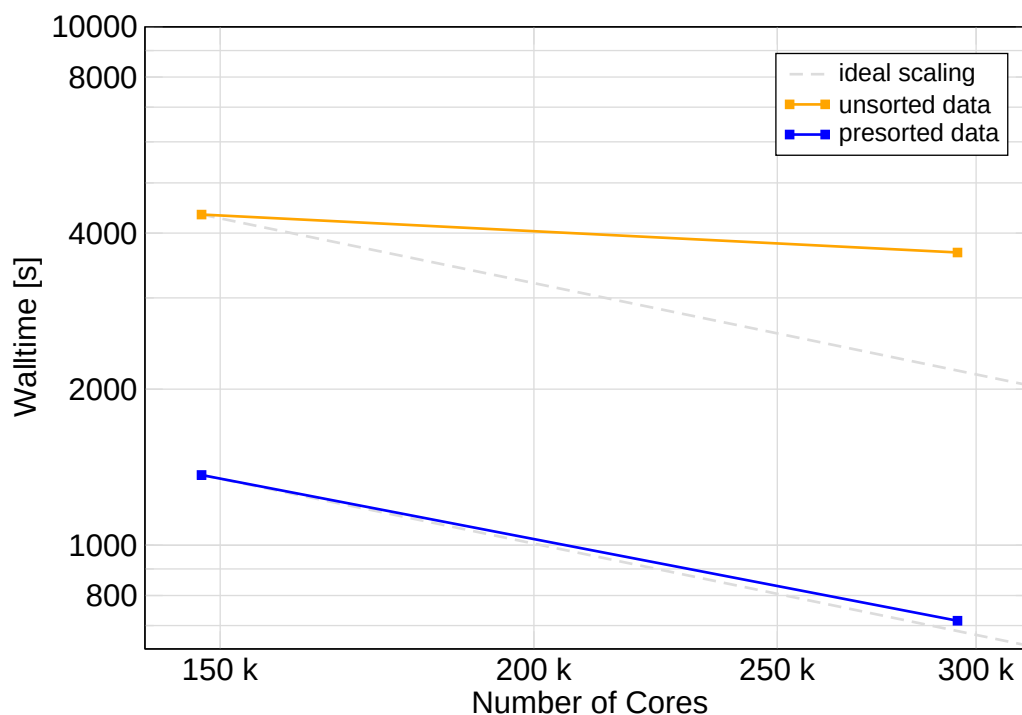
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Passing the Three-Trillion Particle Limit

Open System, 3.011.561.968.121 Particles, $\Delta E = 10^{-2}$, 44.16 Bytes/particle



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Conclusion & Outlook

Take Away Message

- 1 FMMs can offer on-the-fly, single parameter
 - error control
 - runtime minimization
- 2 FMMs can be implemented
 - with a small memory footprint
 - with a small parallel overhead

FMM is part of Scafacos

- FMM implementation is available as part of Scafacos
- FMM will be coupled with GROMACS and LPPOLY